

Making earth science data more accessible: experience with chunking and compression

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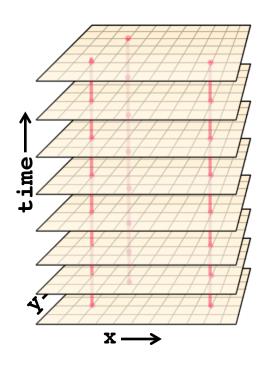
What's the Problem?

Can large multidimensional datasets be organized for fast and flexible

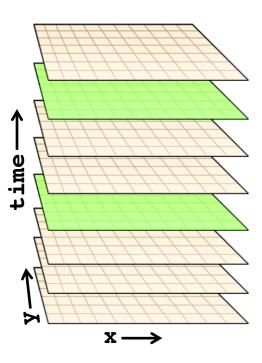
access?

Without multiple versions of the data, what's the best you can do?

Time range access



Spatial access



Conventional storage layout	Time range access	Spatial access
Time varying fastest	Fast	Slow
Time varying slowest	Slow	Fast



Goal is to solve what looks like a little problem that becomes more serious with ...



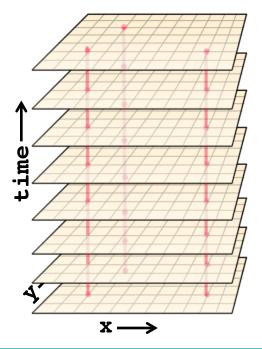
Real data, conventional storage

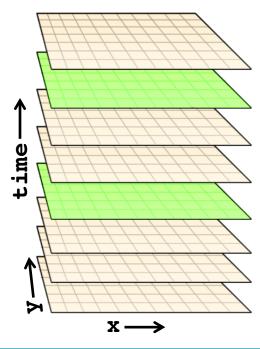
NCEP North American Regional Reanalysis float 200mb_TMP(time=98128, y=277, x=349)

9.5 billion values38 GB of data8.5 GB compressed

Contiguous storage with

- ☐ time varying fastest
- ☐ time varying slowest



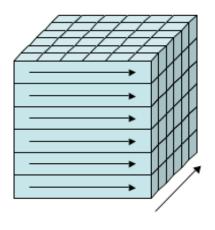


	Time range access	Spatial access
Time varying fastest	0.013 sec	180 sec
Time varying slowest	200 sec	0.012 sec

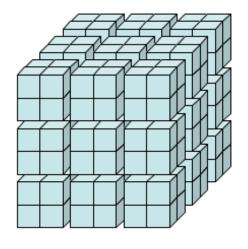
^{*}Single file access time averaged over 100 independent queries, after clearing disk caches. 7200 rpm disk.



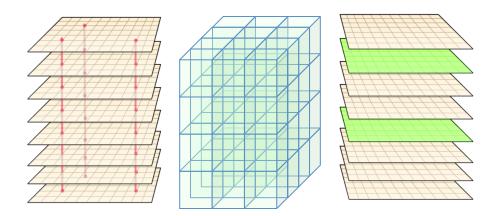
Chunking



index order



chunked



- Storing data in "chunks" along each dimension in a multidimensional variable makes access along any dimension similar
- Slows down fast accesses, but speeds up slow accesses
- How to choose shape and size of chunks for acceptable performance?



Benefits of chunking

- Large performance gains are possible with good choice of chunk shapes and sizes
- Benefits of chunking for compression are underappreciated
- Chunking supports efficiently extending data along multiple dimensions
- So why is use of chunking not widespread?
 - Software to rechunk big datasets is available, but defaults work poorly for some common use cases
 - Specific guidance for how to set chunk shapes is lacking



Importance of chunk shapes

Example: float $200mb_{TMP}(time=98128, y=277, x=349)$

Storage layout, chunk shapes	Read time range (seconds)	Read horizontal slice (seconds)	Performance bias: (slowest / fastest)
Contiguous, for time range	0.013	180	14,000
Contiguous, for spatial slices	200	0.012	17,000
Record chunks, 1 x 277 x 349	690	0.007	98,000
Default chunks, 4673 x 12 x 16	1.4	34	24
36 KB chunks, 92 x 9 x 11	2.4	1.7	1.4
8 KB chunks, 46 x 6 x 8	1.4	1.1	1.2



Chunk shapes

- 2-dimensional analog of chunking is inadequate for common use case of 1D and nD access in an (n+1)-dimensional dataset
- In 2D, want chunks to be same shape as data domain to get same number of chunks in each direction of access
- In 1D and nD access, need to divide chunks read per access equally between 1D and nD domains
- For 3D use case example, balancing 1D and 2D accesses:
 - Let n_{time} = number of time, n_y = number of y, n_x = number of x, N = number of chunks = $(n_{time} \ n_y \ n_x)$ / valuesPerChunk
 - time by y by x chunk shape should be integral, near $n_{time}/N^{\frac{1}{2}}$ by $c n_y/N^{\frac{1}{4}}$ by $1/c n_x/N^{\frac{1}{4}}$ (for any c > 0)
- More detailed guidance will appear soon in <u>Unidata's Developer's Blog</u> and netCDF documentation



Chunking transparency

- Only creators of a dataset need to be concerned with chunk shapes and sizes
- Chunking and compression are *invisible* to reading software, except for performance
- Per-variable rechunking and compression implemented in calls to access libraries
- Per-dimension rechunking and compression supported by nccopy or h5repack utilities



Chunking and compression

- In using netCDF or HDF5 libraries, a chunk is an indivisible unit of disk access, compression, and caching
- In general, smaller chunks mean worse compression
- Smaller chunks improve access times for compressed data, due to less computation for uncompression
- Benchmark times reported above are for uncompressed data
- Including compression doesn't change



Chunk size

faster read access less compression slower to create large chunks

- Chunk size should be at least size of a disk block
- Chunk shape is more important than chunk size for balanced and flexible access in multiple ways
- To re-chunk large datasets, it's best to have lots of memory



Some inadequate chunking advice



19.1 Choosing Chunksizes

How do you pick chunksizes?

- Choosing good chunksizes depends on the access patterns of your data. Are you trying to optimize writing, reading, or both? What are the access patterns at I/O bottlenecks?
- Choose chunksizes so that the subsets of data you are accessing fit into a chunk. That is, the chunks should be as large, or larger than, the subsets you are reading/writing.
- The chunk cache size must also be adjusted for good performance. The cache must be large enough to hold at least one chunk.
- Setting a larger cache (for example, big enough to hold tens or hundreds of chunks) will pay off only if the access patterns support it.
- On today's high-performance systems, large amounts of memory are available (both to the user and as internal hardware caching.) This suggests that chunks and caches should be large, and programs should take large sips of data.

Oops, I wrote that last year, before this example ...



Summary

- By rewriting important datasets using appropriate chunking, you can tailor their access characteristics to make them more useful
- This includes structuring model outputs and remote access observations for flexible and efficient access
- Proper use of chunking can support multiple common query patterns for large datasets
- Specific guidance for how to choose optimal shapes and sizes of multidimensional chunks is currently hard to find
- We will soon be making such advice and software implementing it more widely available



Plans

- Blog on details of what was learned, what remains a mystery
- Improve documentation with better guidance for effective use of chunking and compression
- Improve default chunking in netCDF-4 software
- Improve nccopy implementation of re-chunking to support chunking policies, as in NCO software



For more information:

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unidata.ucar.edu/netcdf/

hdfgroup.org/HDF5/doc/Advanced/Chunking/

Good paper on chunking details:

Optimal chunking of large multidimensional arrays for data warehousing by E. J. Otoo, Doron Rotem, and Sridhar Seshadri, 2008



Thank you!



Benchmark details

- Disk cache in memory cleared before each run
- Measured average clock time to read 100 time ranges and 100 spatial slices
- Each of the 100 time ranges or spatial slices used independent chunks
- Still some speedup from first read to later reads, due to disk caches not in OS memory
- Used local 7200 rpm disk for most tests (44 MB/sec)
- SSD was about 4 x faster in sample comparison runs